## Sparse-matrix algorithms for global eigenvalue problems

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## Global Flow Stability and Control IV

## Why a custom algorithm?

- Treating a discretization of the linearized Navier-Stokes equations as a full matrix is too costly to be practical.
- Most work on global stability has used inverse iteration (for a single mode) or ARPACK library with shift-and-invert (for the dominant few modes)
- The inversion part of both these methods is expensive and generally iterative.


## Why a custom algorithm?

- Treating a discretization of the linearized Navier-Stokes equations as a full matrix is too costly to be practical.
- Most work on global stability has used inverse iteration (for a single mode) or ARPACK library with shift-and-invert (for the dominant few modes)
- The inversion part of both these methods is expensive and generally iterative.
- It is a waste to iterate to convergence something that is in fact a stage of another, outer iteration. Can we modify the eigenvalue algorithm so that a single step of the inversion procedure can be done per iteration?


## From the beginning: direct and inverse iteration

## Direct iteration

$$
\mathbf{x}^{(n+1)}=(\mathbf{A}-s) \mathbf{x}_{n}
$$

$$
\sigma^{(n+1)}=\frac{\mathbf{p} \cdot \mathbf{x}^{(n+1)}}{\mathbf{p} \cdot \mathbf{x}^{(n)}}+s \quad\left(\mathbf{p}=\text { projector; e.g. } \mathbf{p}=\mathbf{x}^{(n) *}\right)
$$

converges to the eigenvalue (if unique) farthest from the shift $s$.

## Inverse iteration

$$
\mathbf{x}^{(n+1)}=(\mathbf{A}-s)^{-1} \mathbf{x}^{(n)}
$$

$$
\sigma^{(n+1)}=\frac{\mathbf{p} \cdot \mathbf{x}^{(n)}}{\mathbf{p} \cdot \mathbf{x}^{(n+1)}}+s \quad\left(p=\text { projector; e.g. } \mathbf{p}=\mathbf{x}^{(n+1) *}\right)
$$

converges to the eigenvalue (if unique) closest to the shift $s$.
Converges quadratically if we let $s=\sigma_{n}$.
Requires matrix inversion.

## Connection with explicit and implicit time integration

## Explicit Euler

$$
\delta t^{-1} \mathbf{x}^{(n+1)}=\left(\mathbf{A}+\delta t^{-1}\right) \mathbf{x}^{(n)}
$$

same as direct iteration with shift $s=-1 / \delta t$.

## Implicit Euler

$$
\mathbf{x}^{(n+1)}=\left(\mathbf{A}-\delta t^{-1}\right)^{-1} \delta t^{-1} \mathbf{x}^{(n)}
$$

same as inverse iteration with shift $s=+1 / \delta t$.

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- As many realized in the past, an already available time-integration algorithm can be used as the core of an eigenvalue iteration.
- As perhaps not as many realized, first-order time integration gives exact eigenvectors and eigenvalues.


## A discretized differential equation is not just any matrix



## Subspace iteration

Using a set of $K$ basis vectors $\mathbf{x}_{k}$ :
orthogonalize $\mathbf{x}_{k}^{(n)}$;
do $\mathbf{x}_{k}^{(n+1)}=(\mathbf{A}-s) \mathbf{x}_{k}^{(n)}$ for $k=1$ to $K$;
let $\sigma_{h k}^{(n+1)}=\left(\mathbf{p}_{h} \cdot \mathbf{x}_{k}^{(n)}\right)^{-1}\left(\mathbf{p}_{h} \cdot \mathbf{x}_{k}^{(n+1)}\right)+s$.
Notes:

- The straightforward choice $\mathbf{p}_{h}=\mathbf{x}_{h}^{(n) *}$ obviates the need to divide by $\mathbf{x}_{h}^{(n) *} \cdot \mathbf{x}_{k}^{(n)}=\delta_{h k}$.
- The eigenvalues of $\sigma_{h k}$ can be extracted by a standard full-matrix library and converge to the eigenvalues of $\mathbf{A}$.
- The rate of convergence is dictated by the first neglected eigenvalue. The leading eigenvectors will converge even if their corresponding eigenvalues are close to each other or multiple.


## Inverse subspace iteration

orthogonalize $\mathbf{x}_{k}^{(n)}$;
do $\mathbf{x}_{k}^{(n+1)}=(\mathbf{A}-s)^{-1} \mathbf{x}_{k}^{(n)}$ for $k=1$ to $K$;
$\sigma_{h k}^{(n+1)}=\left(\mathbf{p}_{h} \cdot \mathbf{x}_{k}^{(n+1)}\right)^{-1}\left(\mathbf{p}_{h} \cdot \mathbf{x}_{k}^{(n)}\right)+s$.

- The standard choice $\mathbf{p}_{h}=\mathbf{x}_{h}^{(n) *}$ this time eliminates the numerator. By personal experience, to try to eliminate the denominator is a bad idea (I could not find this warning in any book).
- K matrix inversions are required. However,
- they can be performed in parallel.


## Arnoldi

A smart way to organize subspace iteration so that the equation

$$
\mathbf{x}_{k}^{(n+1)}=(\mathbf{A}-s) \mathbf{x}_{k}^{(n)}
$$

is already satisfied for $k=1 . . K-1$ and need only be imposed for $k=K$.

- only one application of matrix $\mathbf{A}$ is needed per step instead of $K$, but
- $K$ increases by 1 at every step. At a preset value $K_{\text {max }}$ the algorithm must be restarted, by rotating the current approximations of the leading eigenvectors into the first $K_{\text {min }}$ basis vectors.


## Implicitly restarted Arnoldi

A smart way (incomplete QR) to rotate the basis vectors so as not to ruin the Arnoldi consistence for $k=1$.. $K_{\text {min }}-1$

## Our method (I): approximate-inverse subspace iteration

Orthogonalize $\mathbf{x}_{k}^{(n)}$ and rotate $\sigma_{h k}$ accordingly; construct $\tilde{\mathbf{x}}_{h}^{(n+1)}=\sigma_{h k}^{(n)} \mathbf{x}_{k}^{(n)}$; with $\mathbf{B}$ being an approximate inverse of $(\mathbf{A}-s)$ do

$$
\mathbf{x}_{k}^{(n+1)}=\tilde{\mathbf{x}}_{k}^{(n+1)}+\mathbf{B}\left[\mathbf{x}_{k}^{(n)}-(\mathbf{A}-s) \tilde{\mathbf{x}}_{k}^{(n+1)}\right] \text { for } k=1 \text { to } K ;
$$

$$
\sigma_{h k}^{(n+1)}=\left(\mathbf{p}_{h} \cdot \mathbf{x}_{k}^{(n)}\right)^{-1}\left(\mathbf{p}_{h} \cdot \mathbf{x}_{k}^{(n+1)}\right)
$$

(1) A multigrid algorithm (previously developed for steady-state iteration) provides B.
(2) Depending on the accuracy of the approximation, only part of the eigenvectors are usually found to converge.
(3) A simultaneous iteration of the base flow achieves convergence in unstable cases.

## Our method (II): towards inversion-free Arnoldi (or Jacobi-Davidson?)

Instead of updating each basis vector, increase $K$ by 1 and use the correction as a new basis vector; when $K=K_{\text {max }}$, deflate the search space to $K_{\text {min }}$ by the same incomplete QR procedure as in IRAM.

- If B was an exact inverse, only the last vector would generate a correction and the method would essentially coincide with IRAM.
- Only the vectors with the largest residuals need to generate corrections. Different strategies can be devised and are being experimented with.


## The lid-driven cavity test case

- $2 \frac{1}{2} \mathrm{D}$ stability problem: suboptimal wavenumber (Ramanan \& Homsy 1994).
- $2 \frac{1}{2} \mathrm{D}$ stability problem: $\operatorname{Re}_{\text {cr }}=786 ; \beta=15.8$ (Albensöder \& Kuhlmann 2001).
- 2D stability problem: Re $_{c r}=8018$ (Parolini, Auteri \& Quartapelle 2002).
- 3D stability problem: AFAIK only solved for aspect ratio 1:1:6 (Albensöder \& Kuhlmann 2001)
- 3D stability problem in a 1:1:1 cubical cavity: $\mathrm{Re}_{\text {cr }}$ observed to lie between 2000 and 3000 from direct simulations (Iwahatsu, Ishii \& Kawanura 1989).
- Eigenvalues to follow next...


## Driven cubic box base flow ( $\mathrm{Re}=2000,64^{3}$ ).



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## Driven cubic box spectrum $\left(64^{3}\right)$.



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Driven cubic box spectrum ( $128^{3}$ ).


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Resolution comparison. Re=2000


Resolution comparison. $\mathrm{Re}=2100$


Resolution comparison. $\mathrm{Re}=2200$


Resolution comparison. $\mathrm{Re}=2250$


Resolution comparison. $\mathrm{Re}=2300$


Resolution comparison. $\mathrm{Re}=2400$


High-frequency odd mode $\left(\operatorname{Re}=2000,64^{3}\right)$.


High-frequency even mode ( $\mathrm{Re}=2000,64^{3}$ ).


Low-frequency odd mode $\left(\mathrm{Re}=2000,64^{3}\right)$.


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## Low-frequency even mode ( $\mathrm{Re}=2000,64^{3}$ ).



Zero-frequency least stable mode $\left(\mathrm{Re}=2000,64^{3}\right)$.


- Method I: approximate-inverse subspace iteration with multigrid.
- Method II: IRAM-like version of the above.
- Eigenvalue spectrum of the 3D 1:1:1 lid-driven cavity. $R e_{c r}=2200, \omega=0.50$

Ongoing developments:

- improved parallelization ( $256^{3}, 24+$ modes),
- combined direct-adjoint iteration (Lanczos),
- immersed boundary, non-uniform grid and application to open flows.

